#### VM-based STM

- Is changing the VM acceptable?
- · Benefits
  - Direct memory access
  - Full control over garbage collector (GC)
  - Full control over bytecode instructions behavior
  - Can manipulate thread's header
  - HTM compatible

#### Implicit transaction

atomic{ A = B; B++; }

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#### Or:

stm.STM.xBegin(); A = B; B++; stm.STM.xCommit();

**Implicit transaction** 

Transaction T; T.begin() do{ A.txWrite(B.txRead()); B.txWrite(B.txRead() + 1); }while(!T.commit());

**Explicit transaction** 

- No special transactional instructions
  - Bytecode instructions have two modes
    - · Transactional
    - · Non-transactional
  - Two new bytecode instructions only
  - One copy of the code
- · Works on all data types
  - Memory access is monitored at the bytecode instructions level

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· Supports external libraries

- · Atomic blocks anywhere in the code
  - Saves program state at transaction start
  - Restores the saved state when transaction aborted
  - Monitors less objects

int c=10;	@Atomic			
c = a + 5;	void method(int c){			
atomic{	c = c / 2;			
c = c / 2;	a = c			
a = c;	}			
}				

- · Memory model
  - Direct memory access
    - · Faster write back
  - Raw memory model
    - One code to handle all cases
    - Moving GC compatible (Absolute address is not used)

Instance field: Object address + field offset Static field: Static memory address + field offset Array element: Array address + element size x element index

	Data Type	Base Object	offset	Value	Size	
Obj1.x	int	Obj1	0	20	4	Raw
Obj1.y	double	Obj1	4	46	4	memory model
Obj2.obj	Object (reference)	Obj2	0	0 (i취dex)	4	

Absolute address

- · Write-set
  - Arrays of Primitive + Open Addressing Hashing



- Metadata in the thread header
  - Faster than Java standard ThreadLocal
- · No GC overhead

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- Manually allocates and recycles memory for transactional metadata
- Directly fix write-set only referenced objects

#### **Related Work**



#### Performance

· Linked List



#### Performance

· Vacation

